

# Gas Network Modelling and Optimal Locations for Control under Uncertainty

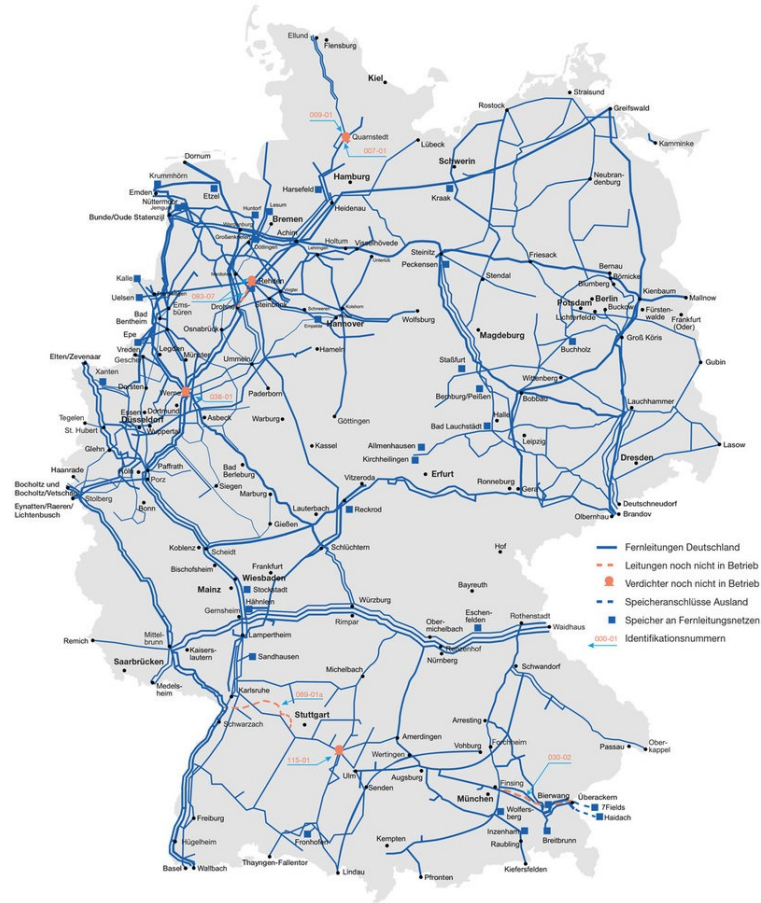
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March 8, Nanzan University

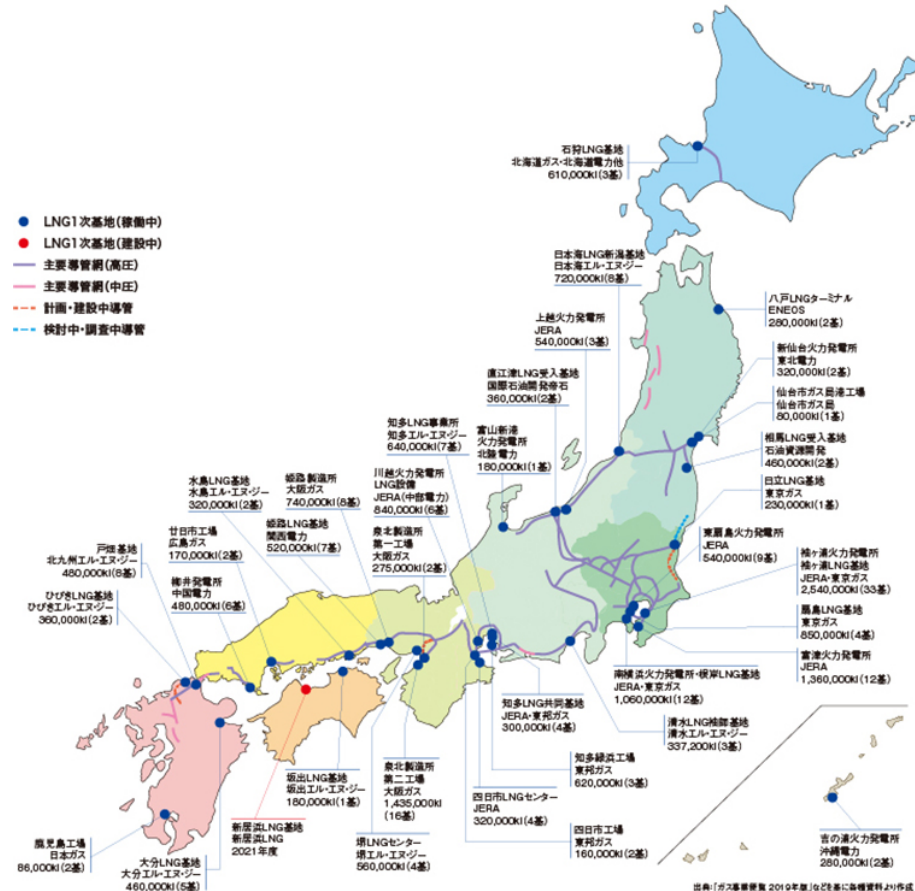
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# Motivation

## Natural Gas Transport



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The isothermal Euler equations for ideal gases:

**(ISO1)**

$$\rho_t + q_x = 0,$$

$$q_t + \left( p + \frac{q^2}{\rho} \right)_x = -\frac{\lambda}{2D} \frac{q|q|}{\rho},$$

**Inlet density & Gas outflow**

$$\rho(t, 0) = \rho_0(t),$$

$$q(t, L) = b(t).$$

**Initial condition**

$$\rho(0, x) = \rho_{ini}(x),$$

$$q(0, x) = q_{ini}(x).$$

see Gugat and Ulbrich (2018): *Lipschitz solutions of initial boundary value problems for balance laws*. Math. Models Methods Appl. Sci., 28(5): 921–951

# Gas Network Modelling

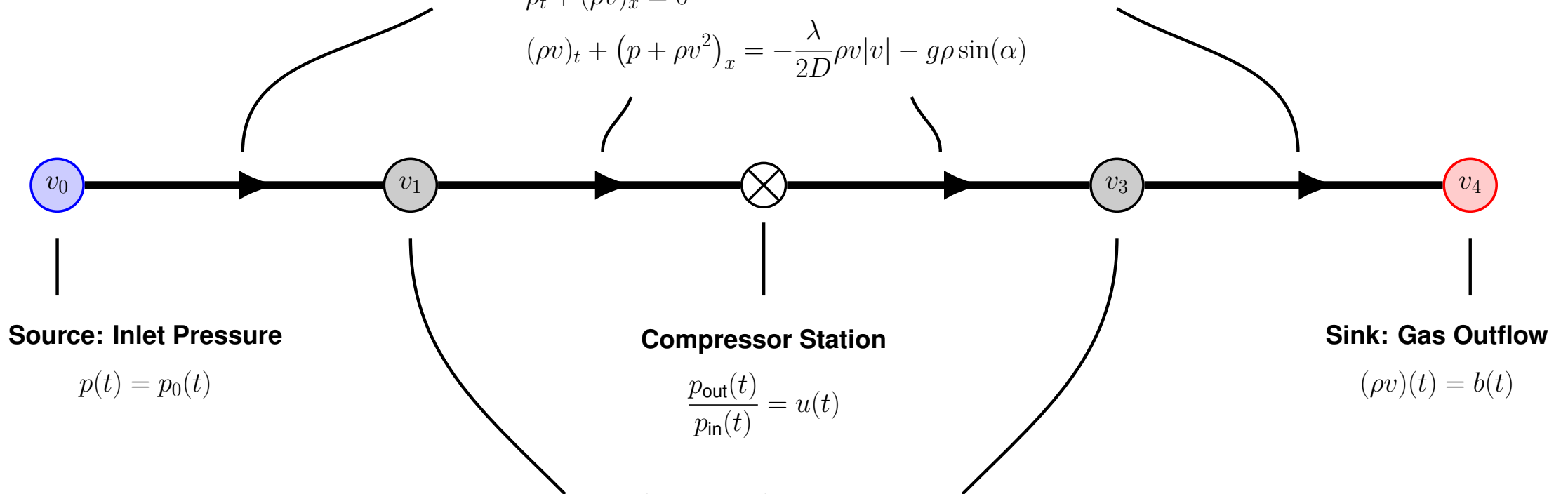
## Gas Flow in Pipeline Networks

$p$	gas pressure	$v$	gas velocity	$g$	gravitational constant
$\rho$	gas density	$\lambda/D$	pipe friction	$\alpha$	pipe slope

### Isothermal Euler Equations

$$\rho_t + (\rho v)_x = 0$$

$$(\rho v)_t + (p + \rho v^2)_x = -\frac{\lambda}{2D} \rho v |v| - g \rho \sin(\alpha)$$



### Coupling Conditions

Conservation of Mass:  $\sum (\rho v)_{in}(t) = \sum (\rho v)_{out}(t),$       Continuity in Pressure:  $p_{in} = p_{out}$



## Model Simplifications

### (ISO1) - quasilinear model

$$\rho_t + q_x = 0$$
$$q_t + \left( p + \frac{q^2}{\rho} \right)_x = -\frac{\lambda^F}{2D} \frac{q|q|}{\rho}$$

### (ISO2) - semilinear model

$$\rho_t + q_x = 0$$
$$q_t + p_x = -\frac{\lambda^F}{2D} \frac{q|q|}{\rho}$$

### (ISO4) - stationary model

$$q_x = 0$$
$$p_x = -\frac{\lambda^F}{2D} \frac{q|q|}{\rho}$$

subsonic flow:  $|v| = \left| \frac{q}{\rho} \right| \ll c$

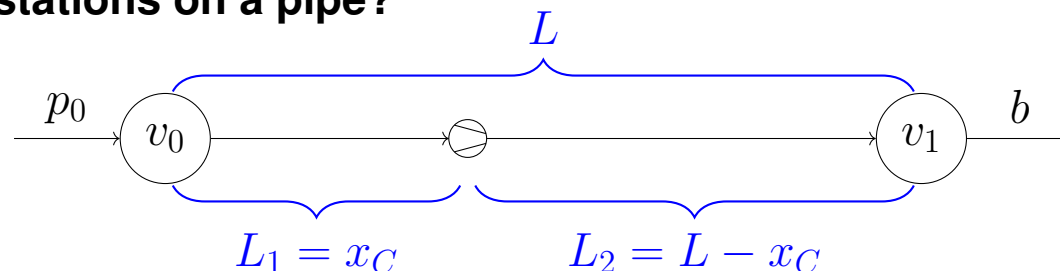
stationary flow:  $\frac{\partial}{\partial t} \equiv 0$

## Mathematical Modelling

- The stationary gas flow for ideal gases on a single pipe is given by

$$q(x) \equiv b \quad (\text{const.}), \quad p(x) = p_0^2 - \phi b |b| x \quad \text{with} \quad \phi = \frac{\lambda}{D} R_S T, \quad x \in [0, L]$$

- Where to place compressor stations on a pipe?



- The stationary gas flow with compressor station for ideal gas is given by

$$p_1^2(x) = p_0^2 - \phi b |b| x \quad x \in [0, L_1]$$

$$p_2^2(x) = u p_0^2 - \phi b |b| (u L_1 + x) \quad x \in [0, L_2]$$

- Consider pressure bounds on the pipe

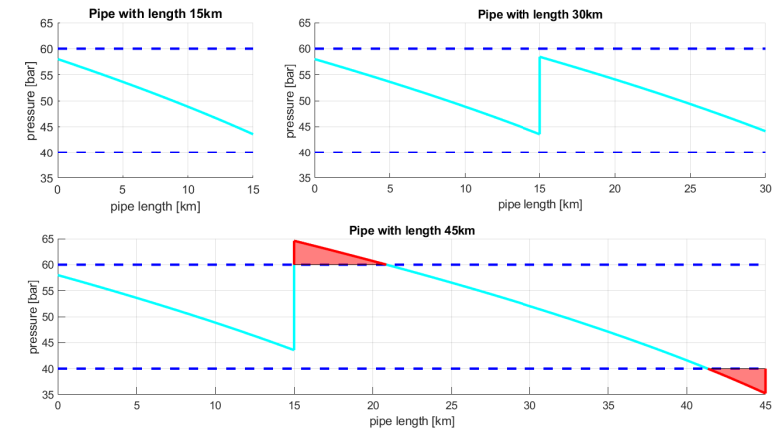
$$\begin{aligned} p_1(x) \in [p_{\min}, p_{\max}] \\ p_2(x) \in [p_{\min}, p_{\max}] \end{aligned} \iff \begin{aligned} p_1(0) \leq p_{\max}, & \quad p_1(L_1) \geq p_{\min}, \\ p_2(0) \leq p_{\max}, & \quad p_2(L_2) \geq p_{\min} \end{aligned}$$

# Optimal Compressor Location

## Deterministic Optimization

Consider the following optimization problem:

$$\text{(OPT 1)} \quad \left\{ \begin{array}{l} \min_{u, x_C} \quad u^2, \\ \text{s.t.} \quad p_1(L_1) \geq p_{\min}, \quad p_2(0) \leq p_{\max}, \quad p_2(L_2) \geq p_{\min}, \\ \quad \quad u \geq 1, \\ \quad \quad x_C \in [0, L]. \end{array} \right.$$

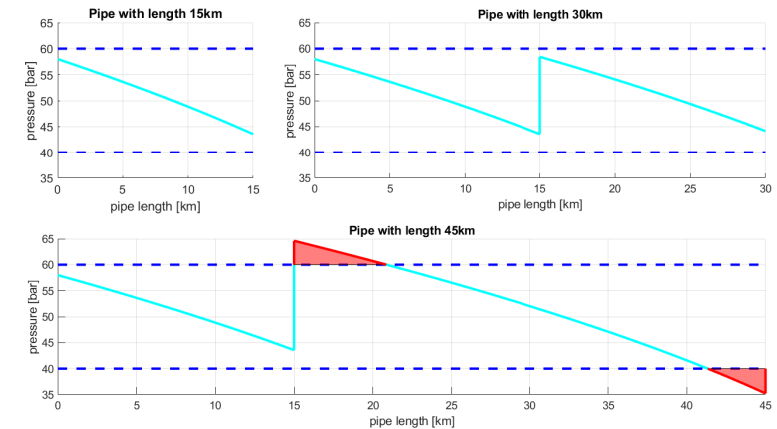


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### Lemma

Let  $p_0 \in [p_{\min}, p_{\max}]$  and  $b > 0$  be given.

- (i) For  $L \leq \frac{p_0^2 - p_{\min}^2}{\phi b |b|}$  every point  $(u, x_C)$  with  $u = 1$  and  $x_C \in [0, L]$  is a solution of the optimization problem (OPT 1).
- (ii) For  $\frac{p_0^2 - p_{\min}^2}{\phi b |b|} < L \leq \frac{p_0^2 + p_{\max}^2 - 2 p_{\min}^2}{\phi b |b|}$  the optimization problem (OPT 1) has a unique solution  $(u^*, x_C^*)$  with  $u^* > 1$  and  $x_C \in [0, L]$ .
- (iii) For  $L > \frac{p_0^2 + p_{\max}^2 - 2 p_{\min}^2}{\phi b |b|}$  the optimization problem (OPT 1) does not have a solution.



# Optimal Compressor Location

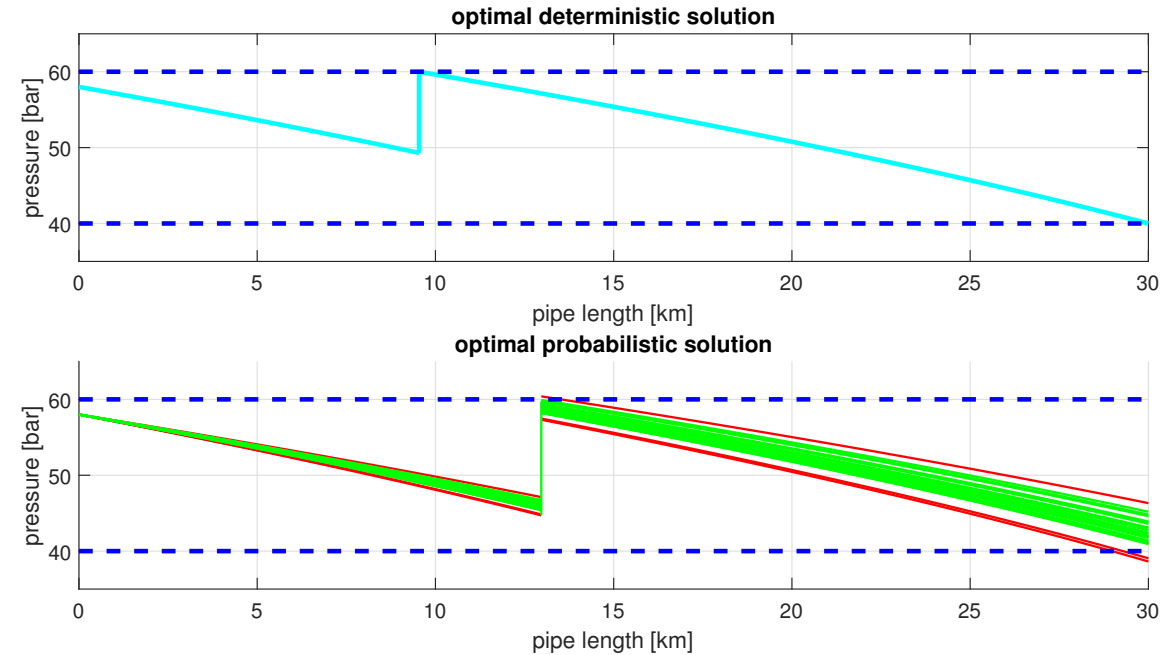
## Probabilistic Optimization

Gas outflow  $b$  is random in the sense that

$$b = \xi(\omega), \quad \xi \sim \mathcal{N}(\mu, \sigma)$$

Consider the following optimization problem:

$$\text{(OPT 2)} \quad \left\{ \begin{array}{l} \min_{u, x_C} u^2, \\ \text{s.t. } \mathbb{P} \left( \begin{array}{l} p_1(L_1) \geq p_{\min} \\ p_2(0) \leq p_{\max} \\ p_2(L_2) \geq p_{\min} \\ p_2(L_2) \leq p_{\max} \end{array} \right) \geq \alpha, \\ u \geq 1, \\ x_C \in [0, L]. \end{array} \right.$$



### Lemma

If  $(u^*, x_C^*)$  with  $u^* > 1$  is a solution of (OPT 2), then the probabilistic constraint is active.

### Theorem

Let  $p_0 \in [p_{\min}, p_{\max}]$  be given.

- (i) If there exists a pair  $(u, x_C)$  with  $u = 1$  and  $x_C \in [0, L]$ , that satisfies the constraints of (OPT 2), then every pair  $(u, x_C)$  with  $u = 1$  and  $x_C \in [0, L]$  is a solution of (OPT 2).
- (ii) If there exist a pair  $(u, x_C)$ , that satisfies the constraints of (OPT 2) and if  $(u, x_C)$  with  $u = 1$  is infeasible for at least one  $x_C \in [0, L]$ , then there exists at least one solution  $(u^*, x_C^*)$  of (OPT 2) with  $u^* > 1$  and  $x_C^* \in [0, L]$ .

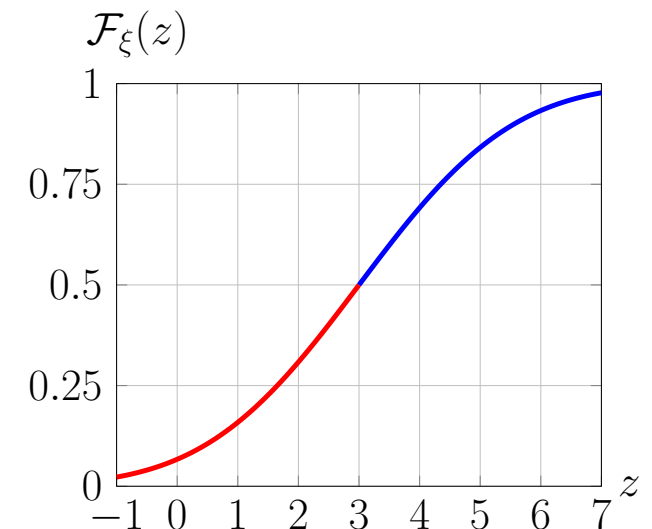
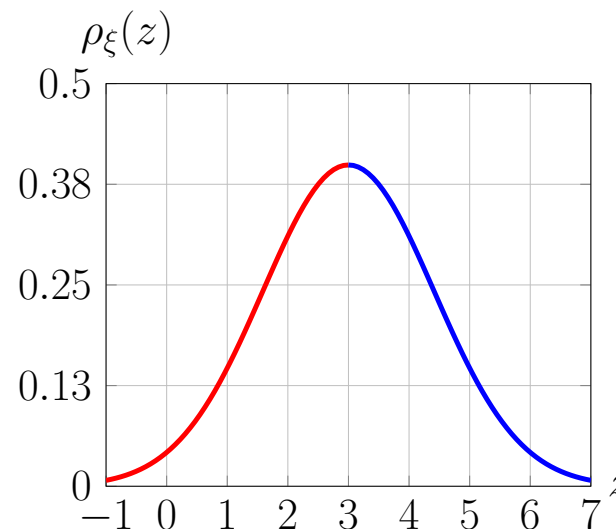
### Theorem

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- (ii) If there exist a pair  $(u, x_C)$ , that satisfies the constraints of (OPT 2) and if  $(u, x_C)$  with  $u = 1$  is infeasible for at least one  $x_C \in [0, L]$ , then there exists at least one solution  $(u^*, x_C^*)$  of (OPT 2) with  $u^* > 1$  and  $x_C^* \in [0, L]$ .

### Theorem

Let  $\alpha > \frac{1}{2}$  be given. For a Gaussian distribution, *Statement (ii) in the last Theorem* guarantees the existence of a unique solution  $(u^*, x_C^*)$  of (OPT 2) with  $u^* > 1$  and  $x_C^* \in [0, L]$ .



# Optimal Compressor Location

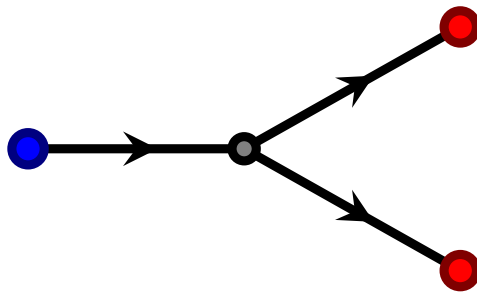
## Mathematical Modelling on Networks

- Consider a connected, directed graph  $\mathcal{G} = (\mathcal{V}, \mathcal{E})$  with vertex set  $\mathcal{V}$  and set of edges  $\mathcal{E}$
- Binary variables  $\delta_i$  states if a compressor location is located on edge  $e_i$
- The stationary gas flow for ideal gas on pipe  $e_i$  is given by

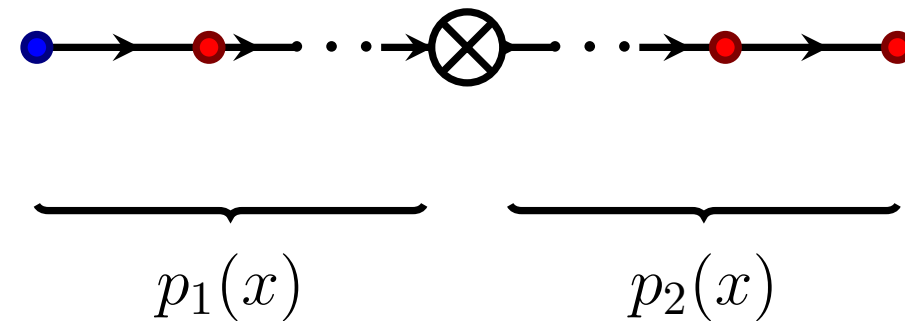
$$p_{i,1}^2(x) = p_{i,1}^2(0) - \phi q_i |q_i| x \quad x \in [0, \delta_i x_{C,i}]$$

$$p_{i,2}^2(x) = (1 - \delta_i + \delta_i u_i) p_{i,1}^2(L_{i,1}) - \phi q_i |q_i| x \quad x \in [0, L - \delta_i x_{C,i}]$$

- Uniqueness in general mainly depends on the graph topology



(a) Scheme of a symmetric graph with one source and two sinks



(b) Scheme of a linear graph with one source and  $n$  sinks



# Optimal Compressor Location

## Deterministic Optimization on Networks

Consider the deterministic optimization problems

$$\text{(OPT 3)} \left\{ \begin{array}{l} \min_{u, x_C, \delta} \|u\|_2^2 = \sum_{i=1}^m u_i^2, \\ \text{s.t. for all } i = 1, \dots, m, \text{ we have} \\ p_{i,1}(L_{i,1}) \geq p_{\min}, \\ p_{i,2}(0) \leq p_{\max}, \\ p_{i,2}(L_{i,2}) \geq p_{\min}, \\ u_i \geq 1, \quad x_{C,i} \in [0, L_i], \quad \delta_i \in \{0, 1\}, \\ \sum_{j=1}^m \delta_j = n_C, \end{array} \right.$$

$$\text{(OPT 4)} \left\{ \begin{array}{l} \min_{u, x_C, \delta, n_C} \|u\|_2^2 = \sum_{i=1}^m u_i^2, \\ \text{s.t. for all } i = 1, \dots, m, \text{ we have} \\ p_{i,1}(L_{i,1}) \geq p_{\min}, \\ p_{i,2}(0) \leq p_{\max}, \\ p_{i,2}(L_{i,2}) \geq p_{\min}, \\ u_i \geq 1, \quad x_{C,i} \in [0, L_i], \quad \delta_i \in \{0, 1\}, \\ \sum_{j=1}^m \delta_j = n_C. \end{array} \right.$$

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### Theorem

For all  $v_i \in \mathcal{V}_{\text{in}}$  let  $p_{i,0} \in [p_{\min}, p_{\max}]$  be given and for all  $v_i \in \mathcal{V}_{\text{out}}$  let  $b_i \geq 0$  be given. Further let a number  $n_C \in \{0, \dots, m\}$  be given.

- (i) If a triple  $(u, x_C, \delta) \in \mathbb{R}^m \times \mathbb{R}^m \times \mathbb{R}^m$  with  $u_j = 1$  for all  $j = 1, \dots, m$  satisfies the constraints of (OPT 3), every triple  $(\mathbf{1}_m, x_C, \delta)$  with  $x_{C,j} \in [0, L_j]$ ,  $\delta_j \in \{0, 1\}$  and  $\sum_{j=1}^m \delta_j = n_C$  is a solution of (OPT 3).
- (ii) If there exists a triple  $(u, x_C, \delta) \in \mathbb{R}^m \times \mathbb{R}^m \times \mathbb{R}^m$  that satisfies the constraints in (OPT 3), and if  $(\mathbf{1}_m, x_C, \delta)$  is infeasible for at least one pair  $(x_C, \delta)$  with  $x_{C,j} \in [0, L_j]$  and  $\delta_j \in \{0, 1\}$ , the optimization problem (OPT 3) has at least one solution.

# Optimal Compressor Location

## Probabilistic Optimization on Networks

Consider the probabilistic optimization problems

$$\text{(OPT 6)} \left\{ \begin{array}{l} \min_{u, x_C, \delta} \|u\|_2^2 = \sum_{i=1}^m u_i^2, \\ \text{s.t. } \mathbb{P} \left( \begin{array}{l} p_{k,1}(L_{k,1}) \geq p_{\min} \\ p_{k,2}(0) \leq p_{\max} \\ p_{k,2}(L_{k,2}) \geq p_{\min} \\ p_{k,2}(L_{k,2}) \leq p_{\max} \end{array} \quad \forall k = 1, \dots, m \right) \geq \alpha, \\ \text{and for all } i = 1, \dots, m, \text{ we have} \\ u_i \geq 1, \quad x_{C,i} \in [0, L_i], \quad \delta_i \in \{0, 1\}, \\ \sum_{j=1}^m \delta_j = n_C, \end{array} \right.$$

$$\text{(OPT 7)} \left\{ \begin{array}{l} \min_{u, x_C, \delta, n_C} \|u\|_2^2 = \sum_{i=1}^m u_i^2, \\ \text{s.t. } \mathbb{P} \left( \begin{array}{l} p_{k,1}(L_{k,1}) \geq p_{\min} \\ p_{k,2}(0) \leq p_{\max} \\ p_{k,2}(L_{k,2}) \geq p_{\min} \\ p_{k,2}(L_{k,2}) \leq p_{\max} \end{array} \quad \forall k = 1, \dots, m \right) \geq \alpha, \\ \text{and for all } i = 1, \dots, m, \text{ we have} \\ u_i \geq 1, \quad x_{C,i} \in [0, L_i], \quad \delta_i \in \{0, 1\}, \\ \sum_{j=1}^m \delta_j = n_C. \end{array} \right.$$

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### Theorem

Let  $p_{i,0} \in [p_{\min}, p_{\max}]$  be given for every node  $v_i \in \mathcal{V}_{\text{in}}$ . Further let a number  $n_C \in \{0, \dots, m\}$  be given.

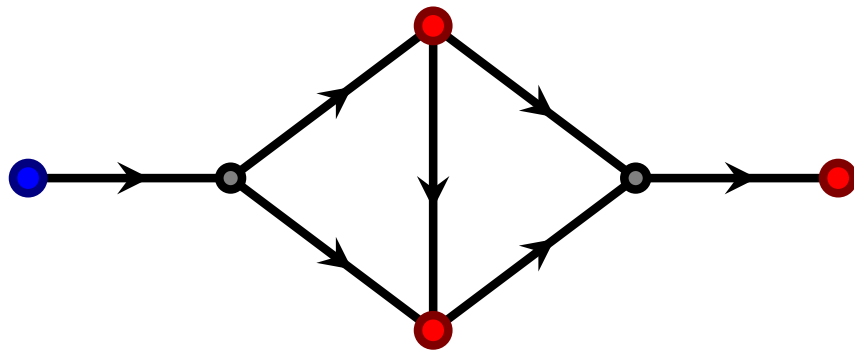
- (i) If there exists a triple  $(u, x_C, \delta) \in \mathbb{R}^m \times \mathbb{R}^m \times \mathbb{R}^m$  with  $u = \mathbf{1}_m$ , that satisfies the constraints of (OPT 6), then every triple  $(\mathbf{1}_m, x_C, \delta)$  with  $x_{C,j} \in [0, L_j]$ ,  $\delta_j \in \{0, 1\}$  and  $\sum_{j=1}^m \delta_j = n_C$  is a solution of (OPT 6).
- (ii) If there exists a triple  $(u, x_C, \delta) \in \mathbb{R}^m \times \mathbb{R}^m \times \mathbb{R}^m$ , that satisfies the constraints of (OPT 6), and if  $(\mathbf{1}_m, x_C, \delta)$  is infeasible for at least one pair  $(x_C, \delta)$  with  $x_{C,j} \in [0, L_j]$  and  $\delta \in \{0, 1\}$ , then (OPT 6) has at least one solution.



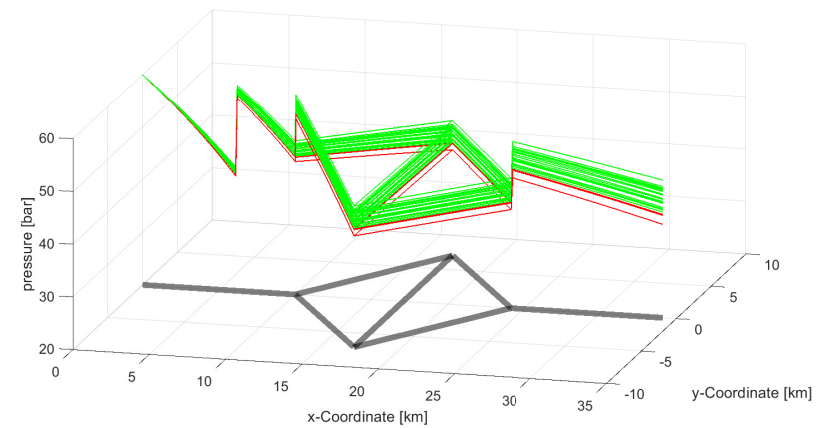
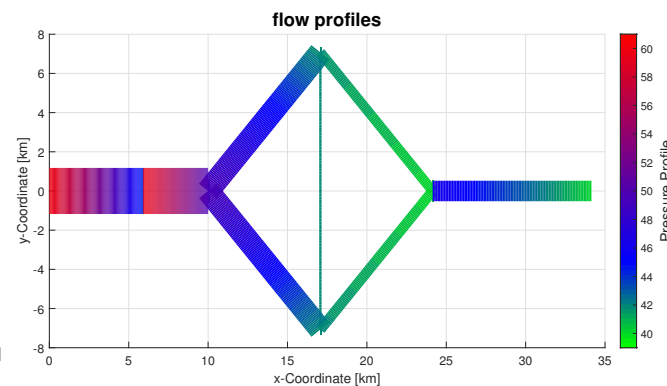
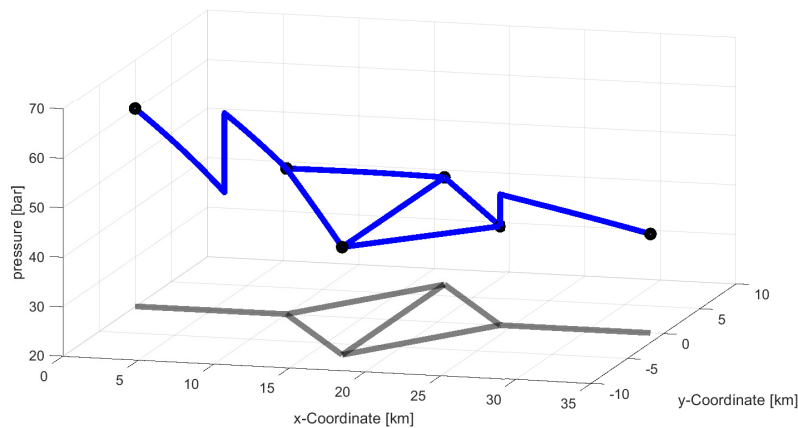
# Optimal Compressor Location

## A Numerical Example on a Diamond Graph

A scheme of a diamond graph with 1 source (blue) and 3 sinks (red):



Variable	Letter	Value	Unit
inlet pressure	$p_0$	60	bar
lower pressure bound	$p_{\min}$	40	bar
upper pressure bound	$p_{\max}$	60	bar
gas outflow (=mean value)	$b (= \mu)$	[90, 60, 120]	kg/m <sup>2</sup> s
covariance matrix	$\Sigma$	diag(2.25, 2.25, 2.25)	
speed of sound in the gas	$c$	343	m/s
pipe friction coefficient	$\lambda^F$	0.1	
pipe diameter	$D$	0.5	m
specific gas constant	$R_S$	515	J/kg K
gas temperature	$T$	293	K
probability level	$\alpha$	0.8	

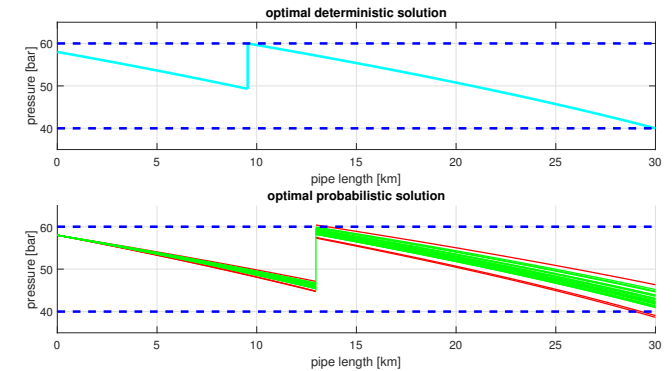


# Optimal Compressor Location

## A Numerical Example for Transient Gas Flow

- Consider the probabilistic example for the stationary flow on a single pipe
- Solve the probabilistic constraint optimization problem for  $\alpha = 90\%$

$$\text{(OPT 2)} \quad \left\{ \begin{array}{l} \min_{u, x_C} u^2, \\ \text{s.t. } \mathbb{P} \left( \begin{array}{l} p_1(L_1) \geq p_{\min} \\ p_2(0) \leq p_{\max} \\ p_2(L_2) \geq p_{\min} \\ p_2(L_2) \leq p_{\max} \end{array} \right) \geq \alpha \\ u \geq 1, \\ x_C \in [0, L]. \end{array} \right.$$

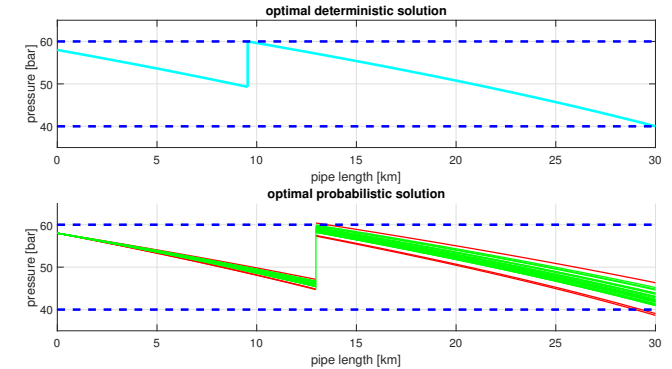


# Optimal Compressor Location

## A Numerical Example for Transient Gas Flow

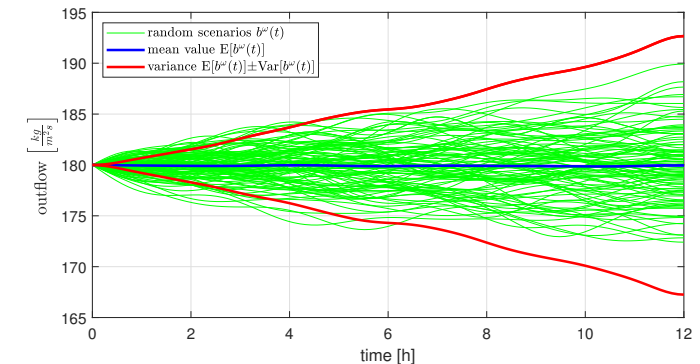
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$$\text{(OPT 2)} \quad \left\{ \begin{array}{l} \min_{u, x_C} u^2, \\ \text{s.t. } \mathbb{P} \left( \begin{array}{l} p_1(L_1) \geq p_{\min} \\ p_2(0) \leq p_{\max} \\ p_2(L_2) \geq p_{\min} \\ p_2(L_2) \leq p_{\max} \end{array} \right) \geq \alpha \\ u \geq 1, \\ x_C \in [0, L]. \end{array} \right.$$



- Randomize the boundary data in time by a *Wiener process*
- The probabilistic robustness of the steady state control is

$$\mathbb{P} \left( \begin{array}{l} p_1(t, L_1) \geq p_{\min} \\ p_2(t, 0) \leq p_{\max} \\ p_2(t, L_2) \geq p_{\min} \\ p_2(t, L_2) \leq p_{\max} \end{array} \quad \forall t \in [0, T] \right) = 85.74\%$$



# Optimal Compressor Location



## A Numerical Example for Transient Gas Flow

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## References

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